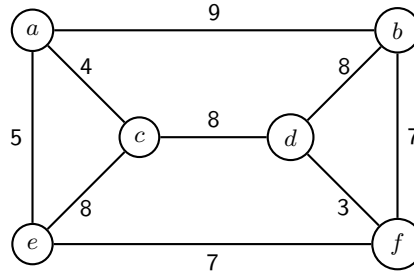


### Problem 1 — Building an MST

Given the weighted graph



Construct an MST.

Iteration	Edge added	Vertex Added
0	—	<i>b</i>
1	<i>bf</i>	<i>f</i>
2	<i>df</i>	<i>d</i>
3	<i>ef</i>	<i>e</i>
4	<i>ae</i>	<i>a</i>
5	<i>ac</i>	<i>c</i>

The tree  $T$ , where  $E(T) = \{bf, df, ef, ae, ac\}$ , is an MST, by Prim's algorithm.

### Problem 2 — Proving a Statement

Prove that if  $G$  is 3-regular and bipartite, it cannot have a bridge.

**Proof:**

Suppose not. Then  $G$  has a bridge  $e = uv$ ,  $e \in E(G)$ . Let  $K$  be the component of  $G - e$  that contains  $u$ . Since  $G$  is bipartite, then  $K$  is bipartite. Also,  $K$  is nearly 3-regular, except for  $u$ , which is 2-regular (since we deleted one of its edges).

Consider the bipartition of  $K$ , call them  $A$  and  $B$ . Since  $u \in K$ ,  $u$  is in one of the bipartitions, let's put it in  $A$ . Observe that the number of edges in  $A$  is the same as in  $B$  (by definition of a bipartition). In  $A$ , there are  $3(|A| - 1) + 2$  edges, and in  $B$ , there are  $3|B|$  edges. However, the number of edges in  $A$  is not divisible by 3 but the number of edges in  $B$  are. This means that both partitions have a different number of edges — a contradiction. Therefore,  $G$  cannot have a bridge.